

Power indices
based on
ordinal games

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Aggregation of ordinal data

and

open problems related to possibilistic
approximation of set functions

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Cooperative games and
Multicriteria decision aid

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1. Introduction

- Set of players $N = \{1, \dots, n\}$
- Game $v: P(N) \rightarrow \mathbb{R}$

Particular cases :

- Probabilistic game p :

$$p(\emptyset) = 0, p(R) = \sum_{i \in R} p(i) \quad \forall R \subseteq N$$

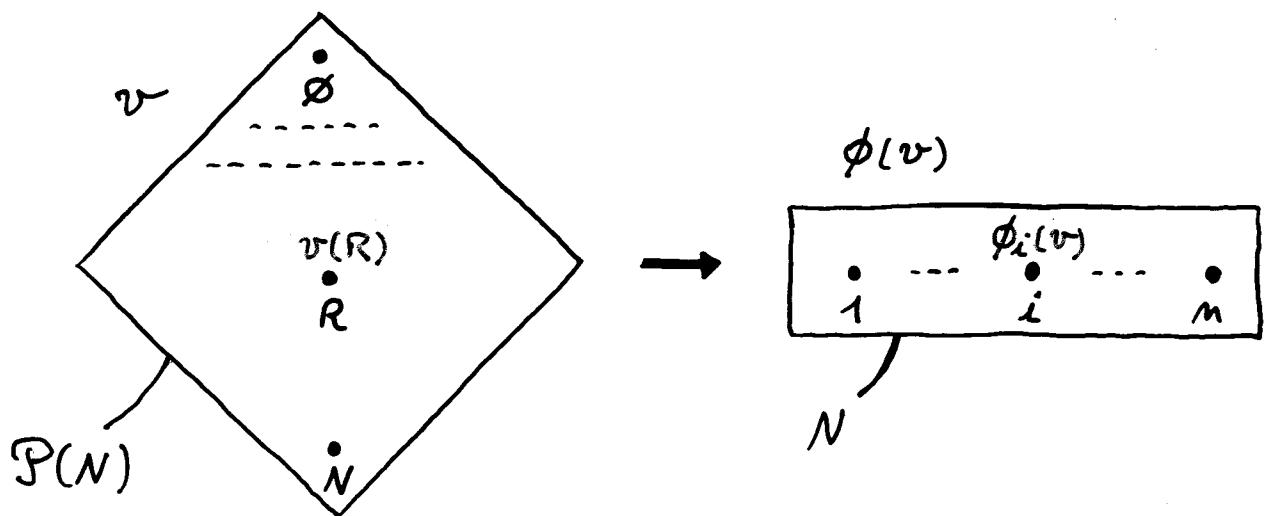
- Possibilistic game π :

$$\pi(\emptyset) = 0, \pi(R) = \vee_{i \in R} \pi(i) \quad \forall R \subseteq N$$

- Power indices $\phi: N \rightarrow \mathbb{R}$

1.1 Cardinal context

Given a game v , we are searching for power indices $\phi(v)$.



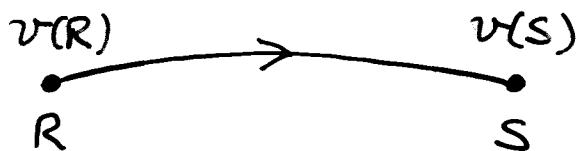
1.2 Ordinal context

$$v' \equiv v \stackrel{\text{def}}{\iff} (v'(R) \geq v'(S) \iff v(R) \geq v(S) \quad \forall R, S \in N)$$

$\Leftrightarrow \exists \psi: R \rightarrow R$ strictly \uparrow s.t. $\alpha' = \psi(\alpha)$

An ordinal game ν is defined by a total preorder V on $\wp(N)$.

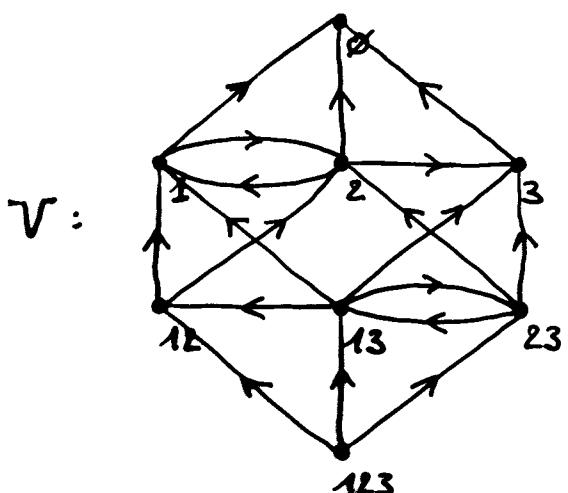
(i.e. a strongly complete and transitive relation \mathcal{V} on $\mathcal{P}(N)$)



Representative Boolean matrix V :

$$V(R, S) = \begin{cases} 1 & \text{if } v(R) \geq v(S) \\ 0 & \text{else} \end{cases}$$

For instance :



Example : Ordinal possibilistic game

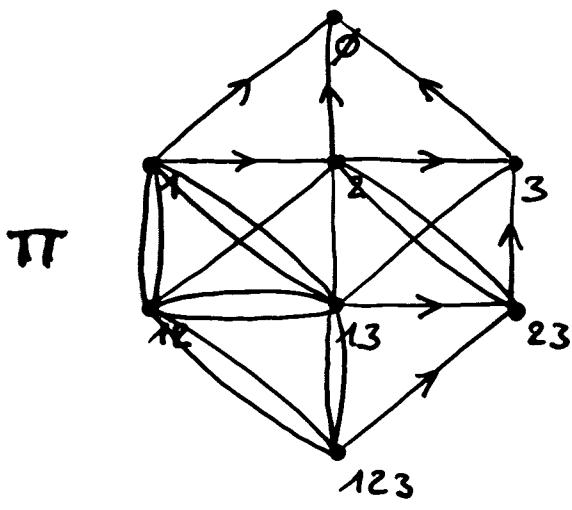
$$\pi(\emptyset) = 0$$

$$\pi(R) = \bigvee_{i \in R} \pi(i)$$

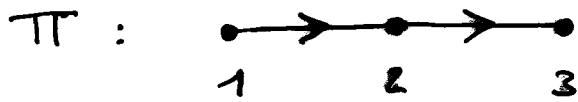


- $\pi(\emptyset, i) = 0 \quad \forall i \in N \quad (\text{i.e. } \pi(i) > \pi(\emptyset))$
- $\pi(R, S) = \begin{cases} 1 & \text{if } \bigvee_{i \in R} \pi(i) \geq \bigvee_{i \in S} \pi(i) \\ 0 & \text{else} \end{cases}$

For instance :



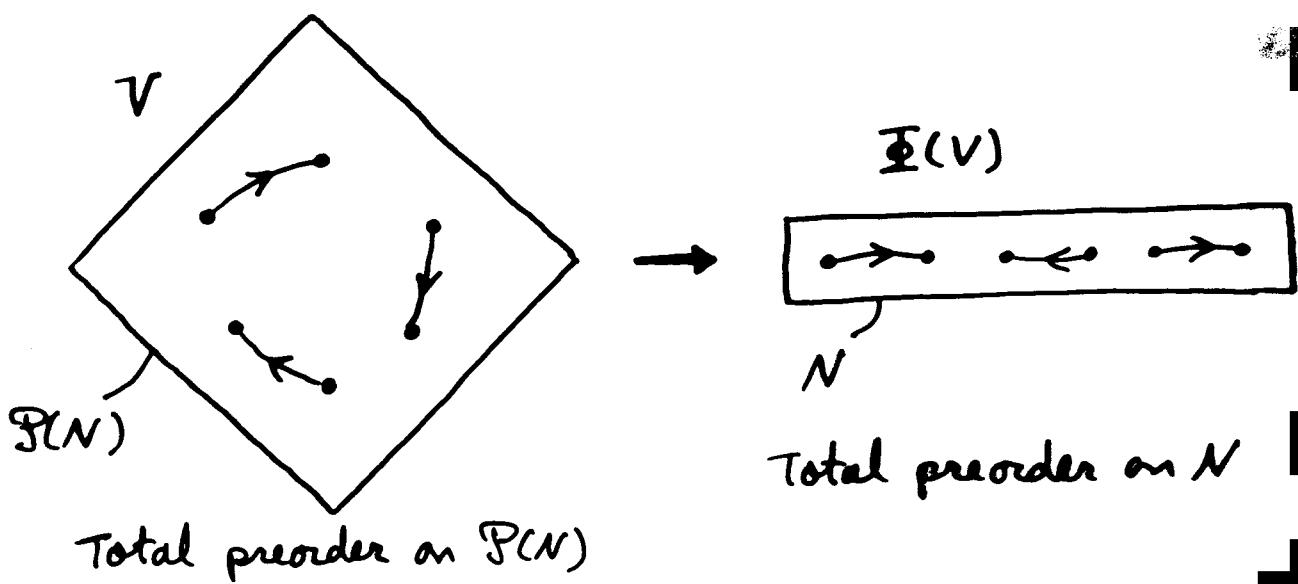
The ordinal possibilistic game π can be assimilated with its induced total preorder on N :



The power indices ϕ are defined by a total preorder Φ on N .

$$\begin{array}{ccc} \phi_i & \xrightarrow{\quad} & \phi_j \\ i & & j \\ \Phi_{ij} = \begin{cases} 1 & \text{if } \phi_i \geq \phi_j \\ 0 & \text{else} \end{cases} \end{array}$$

Given an ordinal game V , we are searching for ordinal power indices $\Phi(V)$:



Numerical representation of Φ :

Let Φ be a binary relation on N . Then Φ is a total preorder if and only if $\exists \phi: N \rightarrow \mathbb{R}$ s.t.

$$\Phi_{ij} = \begin{cases} 1 & \text{if } \phi_i \geq \phi_j \\ 0 & \text{else} \end{cases}$$

ϕ can be defined as $\phi_i = \sum_{j \in N} \Phi_{ij}$

2. Probabilistic approximation

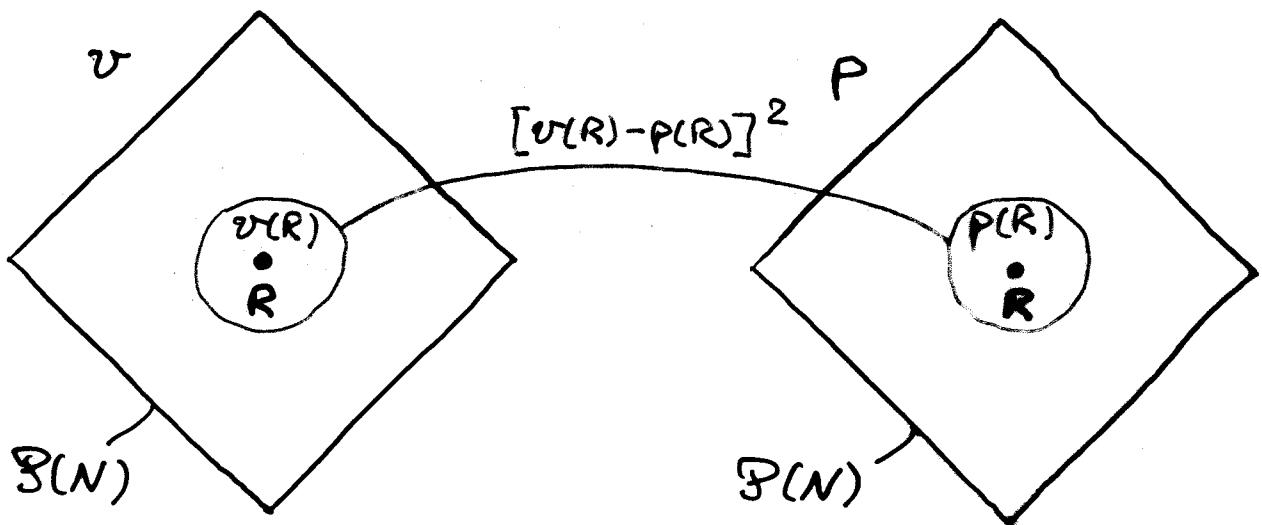
2.1 Cardinal context

The best linear approximation to v , defined as

$$P(R) = P_0 + \sum_{i \in R} P_i$$

corresponds to minimizing

$$d(v, p) = \sum_{R \subseteq N} [v(R) - p(R)]^2$$



It has been characterized by Hammer and Holzman (1992):

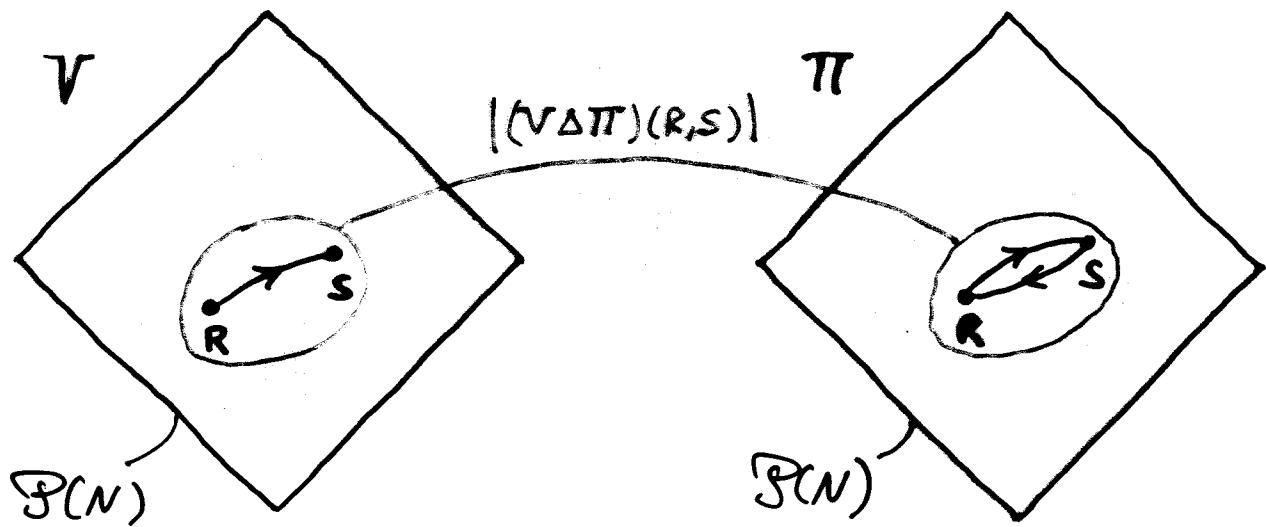
$$\left\{ \begin{array}{l} P_0 = \frac{1}{2^m} \sum_{R \subseteq N} (m - |R| + 1) v(R), \quad |R| = r \\ P_i = b_i(v) = \frac{1}{2^{m-1}} \sum_{R \subseteq N, i \in R} [v(R \setminus i) - v(R)] \quad \forall i \in N \end{array} \right.$$

2.2 Ordinal context

The best probabilistic approximation π to V corresponds to minimizing

$$d(V, \pi) = |V \Delta \pi| = |(V \cup \pi) \setminus (V \cap \pi)|$$

and can be computed with a polynomial time algorithm.

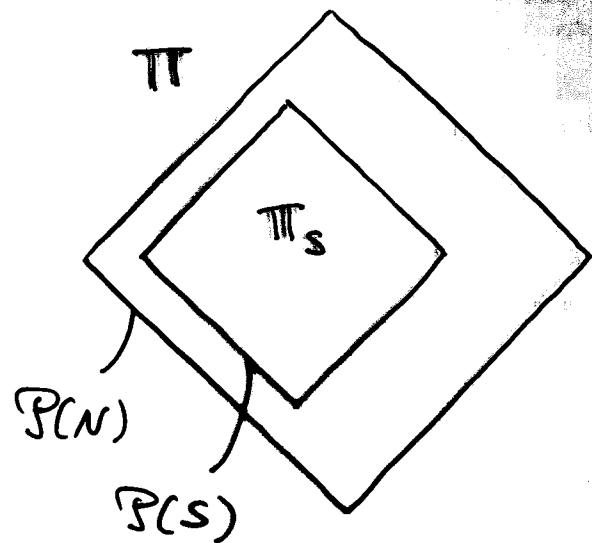
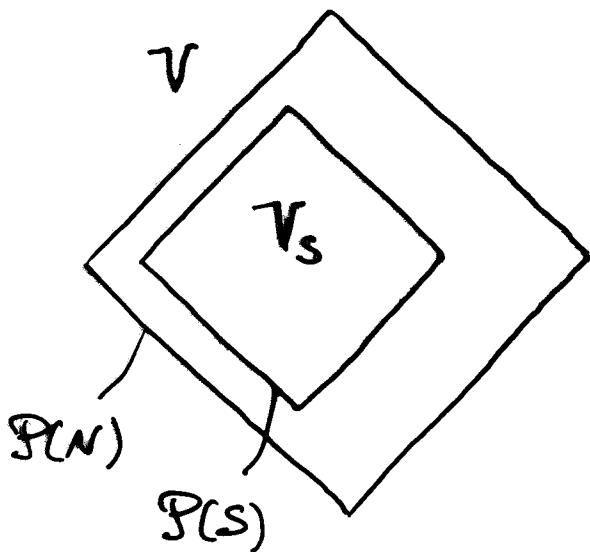


$$|(V \Delta \pi)(R, S)| = \begin{cases} 0 & \text{if } \begin{array}{c} R \xrightarrow{\hspace{1cm}} S \\ R \xrightarrow{\hspace{1cm}} S \end{array} \\ 1 & \text{if } \begin{array}{c} R \xrightarrow{\hspace{1cm}} S \\ R \xleftarrow{\hspace{1cm}} S \end{array} \\ 1 & \text{if } \begin{array}{c} R \xleftarrow{\hspace{1cm}} S \\ R \xrightarrow{\hspace{1cm}} S \end{array} \\ 2 & \text{if } \begin{array}{c} R \xrightarrow{\hspace{1cm}} S \\ R \xleftarrow{\hspace{1cm}} S \end{array} \end{cases}$$

V	π
$R \xrightarrow{\hspace{1cm}} S$	$R \xrightarrow{\hspace{1cm}} S$
$R \xrightarrow{\hspace{1cm}} S$	$R \xleftarrow{\hspace{1cm}} S$
$R \xleftarrow{\hspace{1cm}} S$	$R \xrightarrow{\hspace{1cm}} S$
$R \xrightarrow{\hspace{1cm}} S$	$R \xleftarrow{\hspace{1cm}} S$

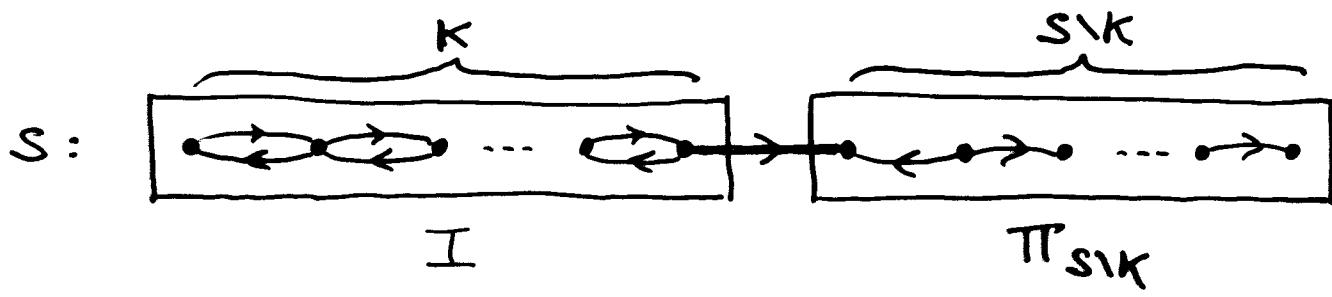
Some notations :

Let $S \subseteq N$. We set :



- $V_S :=$ restriction of V to S
- $z_S^* := \min_{\Pi_S} |V_S \Delta \Pi_S|$
- $\Pi_S^* := \arg \min_{\Pi_S} |V_S \Delta \Pi_S|$
- $\forall K \subseteq S, [I_K, \Pi_{S \setminus K}]$ denotes the total preorder Π_S on S such that

$$\begin{cases} \Pi_S(i,j) = 1 & \text{if } i,j \in K \\ \Pi_S(i,j) = 0 & \text{if } i \in S \setminus K, j \in K \\ \Pi_S = \Pi_{S \setminus K} & \text{on } S \setminus K \end{cases}$$



The result (Crama, Marichal, Mathonet)

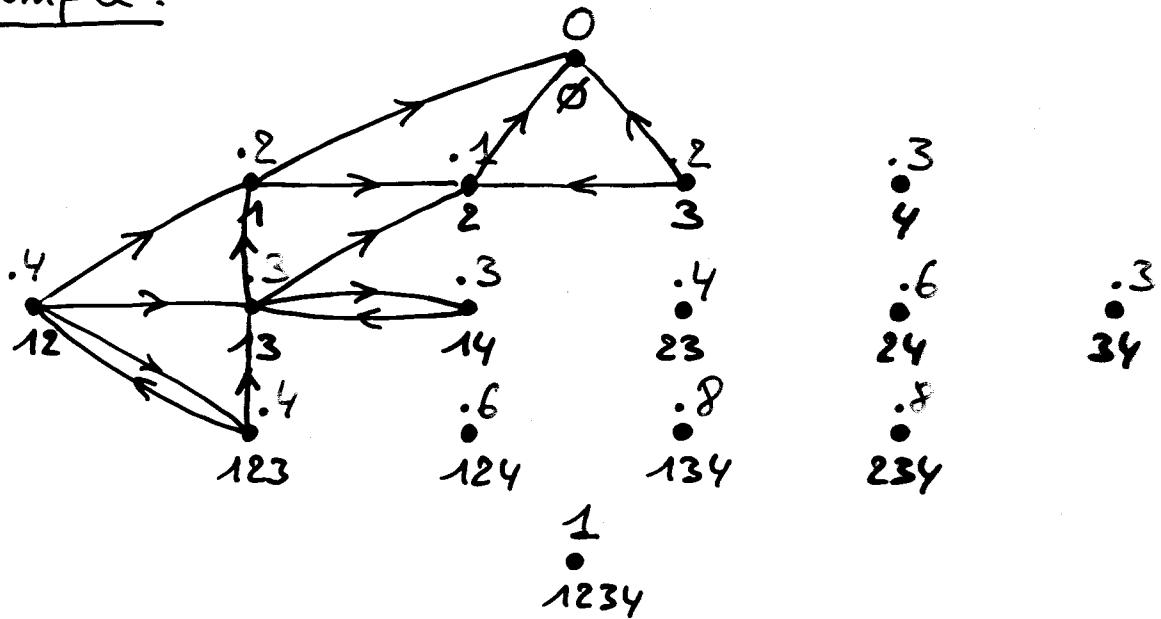
$$z_s^* = \min_{K \subseteq S} |V_S \Delta [I_K, \pi_{s|K}^*]| \quad \forall s \subseteq N$$

Let $S \subseteq N$ and assume that z_T^*, π_T^* are known $\forall T \subseteq S, T \neq S$.

\Rightarrow One can compute z_s^*, π_s^* in $O(2^m \cdot 4^m) = O(8^m)$.

\Rightarrow One can compute $\pi_\emptyset^*, \pi_i^*, \pi_{ij}^*, \dots, \pi_N^*$ in $O(2^m \cdot 8^m) = O(16^m) = O(k^4)$ where $k = 2^m = |\mathcal{P}(N)|$.

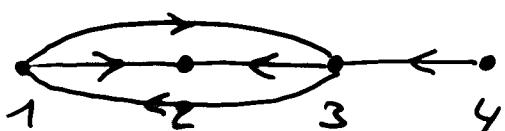
Example:



Optimal solution (unique):

$$\pi(4) > \pi(3) = \pi(1) > \pi(2)$$

$\pi :$



3 Axiomatic approach

3.1 Cardinal context

Given v , one is searching for $\phi(v)$.

Axioms

- Linearity (L) : $\begin{cases} \phi(v+w) = \phi(v) + \phi(w) \\ \phi(c \cdot v) = c \phi(v) \end{cases}$
- Dummy (D) : $v(Ri) = v(R) + \alpha(i) \quad \forall R \subseteq N \setminus i \Rightarrow \phi_i(v) = v(i)$
- Monotonicity (M) : v monotonic $\Rightarrow \phi_i(v) \geq 0 \quad \forall i \in N$
- Symmetry (S) : $\forall \sigma = \text{permutation on } N, \phi_i(v) = \phi_{\sigma(i)}(\sigma v) \quad \forall i \in N$

Weber (1988) :

If $\phi(v)$ fulfills L, D, M, S, then

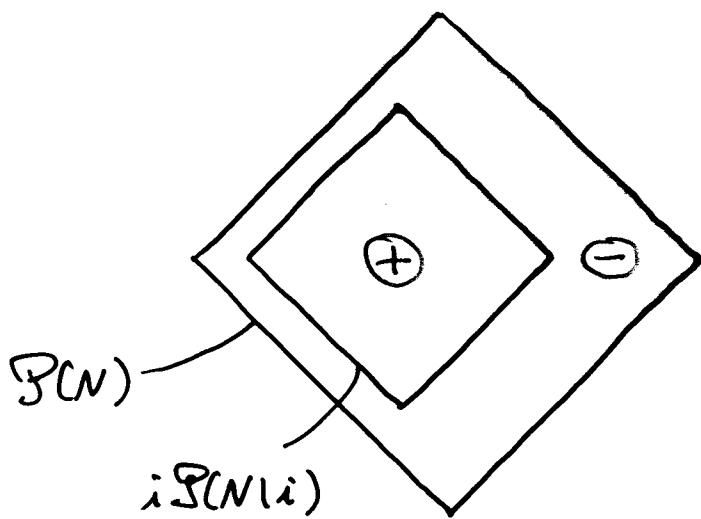
$$\phi_i(v) = \sum_{R \subseteq N \setminus i} p_r [v(Ri) - v(R)] \quad \forall i \in N$$

with $p_r \geq 0, \sum_{R \subseteq N \setminus i} p_r = 1, r = |R|$

For instance :

- Banzhaf power indices:

$$\begin{aligned}
 b_i(v) &= \frac{1}{2^{m-1}} \sum_{R \subseteq N \setminus i} [v(R_i) - v(R)] \\
 &= \frac{1}{2^{m-1}} \left[\underbrace{\sum_{R \subseteq N \setminus i} v(R_i)}_{\oplus} - \underbrace{\sum_{R \subseteq N \setminus i} v(R)}_{\ominus} \right]
 \end{aligned}$$



- Shapley power indices

$$s_i(v) = \frac{1}{m!} \sum_{R \subseteq N \setminus i} r! (m-r-1)! [v(R_i) - v(R)]$$

3.2 Ordinal context

Given V , one is searching for $\Phi(V)$.

(V is monotonic)

Axioms

- $\Phi(V)$ = total preorder

- strong completeness:

$$\forall i, j \in N : \Phi_{ij}(V) = 1 \text{ or } \Phi_{ji}(V) = 1$$

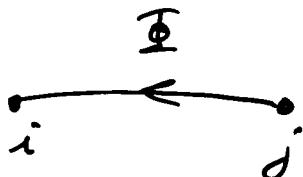
- transitivity:

$$\forall i, j, k \in N : \Phi_{ij}(V) = 1, \Phi_{jk}(V) = 1 \Rightarrow \Phi_{ik}(V) = 1$$

- Null player (N): $(v(R_i) = v(R) \quad \forall R \subseteq N \setminus i)$

$$V(R_i, R) = V(R, R_i) = 1 \quad \forall R \subseteq N \setminus i$$

$$\Rightarrow \Phi_{ji}(V) = 1 \quad \forall j \in N \setminus i$$



- Symmetry (S):

$\forall \sigma = \text{permutation on } N,$

$$\Phi_{ij}(V) = \Phi_{\sigma(i)\sigma(j)}(\sigma V) \quad \forall i, j \in N$$

- Increasingness (I):

Given V , let $V' = V$ except that
 $\exists R \subseteq N \setminus \{i, j\}$ s.t. $v'(R_i) > v(R_i)$ or $v'(R_j) < v(R_j)$.

Then $\Phi_{ij}(V') \geq \Phi_{ij}(V)$ and $\Phi_{ji}(V') \leq \Phi_{ji}(V)$

- Dominance (D):

$$V(R_i, R_j) = 1 \quad \forall R \subseteq N \setminus \{i, j\} \text{ (i.e. } v(R_i) \geq v(R_j)) \Rightarrow \Phi_{ij}(V) = 1$$

- Strong dominance (SD):

$$V(R_i, R_j) = 1 \quad \forall R \subseteq N \setminus \{i, j\} \quad \left. \begin{array}{l} \\ \exists R': V(R'_j, R'_i) = 0 \end{array} \right\} \Rightarrow \Phi_{ji}(V) = 0$$

- Independence of irrelevant coalitions (IC):

$\Phi_{ij}(V)$ does not depend on pairs $(R, S) \subseteq N \times N$ for which
 $i, j \notin R \cup S$ or $i, j \in R \cap S$ or $i, j \in R \setminus S$ or $i, j \in S \setminus R$

Example (Marichal, Mathonet)

$$\phi_i(V) = \sum_{R, S \subseteq N \setminus \{i\}} P_R^S [V(R_i, S) - V(S, R_i)]$$

$$= \sum_{R, S \subseteq N \setminus \{i\}} P_R^S \operatorname{sgn}[v(R_i) - v(S)]$$

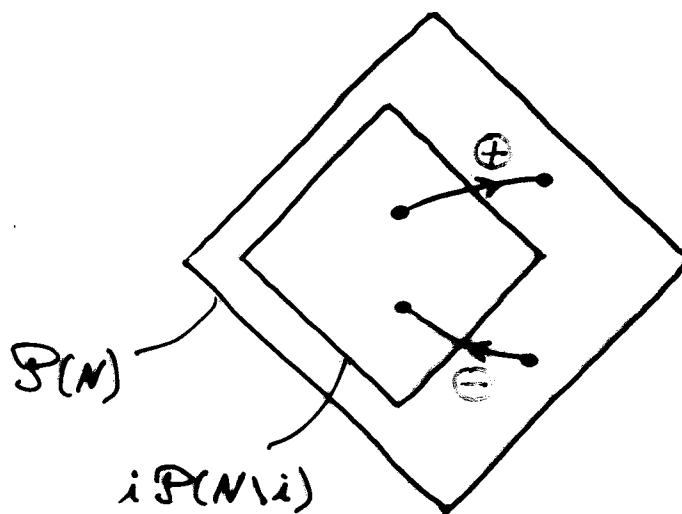
$$(r = |R|, s = |S|) \quad P_R^S > 0$$

fulfills N, S, I, D, SD, IC .

For instance :

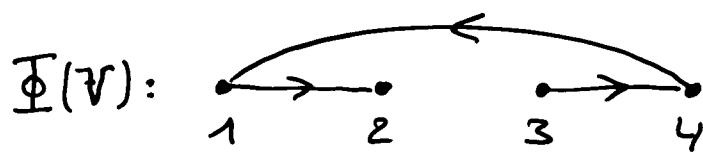
- Ordinal Banzhaf power indices

$$b_i(v) = \sum_{R, S \subseteq N \setminus i} \operatorname{sgn}[v(R_i) - v(S)]$$



Back to the example ($\pi(4) > \pi(3) = \pi(1) > \pi(2)$) :

$$b_3(v) > b_4(v) > b_1(v) > b_2(v)$$



- Ordinal Shapley power indices :

$$s_i(v) = \sum_{R, S \subseteq N \setminus i} r! s! (m-r-1)! (m-s-1)! \operatorname{sgn}[v(R_i) - v(S)]$$